Viewing Graph Solvability in Structure from Motion

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ISPRS TC II online talk series – September 20, 2023

Outline

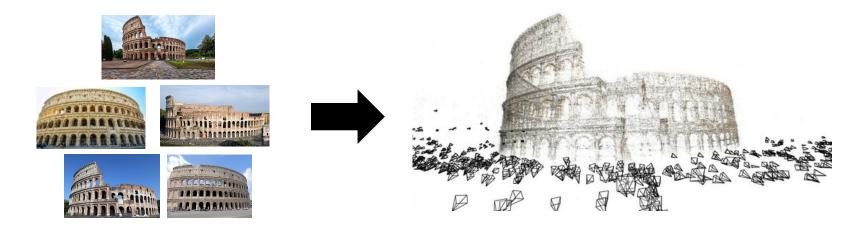
- Introduction
- Calibrated Case
- Uncalibrated Case
- Calibrated vs Uncalibrated
- Conclusion

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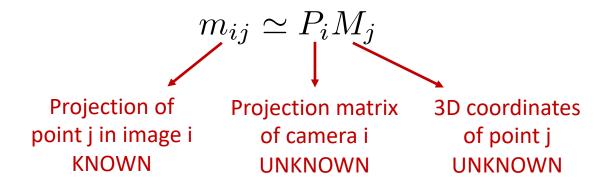
The goal of **structure from motion** (SfM) is to recover both camera motion and scene structure, starting from point correspondences in multiple images:

- camera motion = camera matrices/poses;
- scene structure = 3D coordinates of points.



O. Ozyesil, V. Voroninski, R. Basri, A. Singer. A survey of structure from motion. Acta Numerica (2017).

Formally, the task is to compute **camera matrices** P_i and **coordinates of 3D points** M_i starting from image points m_{ij} such that the following equation is best satisfied:



In the calibrated case, calibration matrices are known and projection matrices consist of **rotations** and **translations**: $P_i = K_i[R_i \ \mathbf{t}_i]$

Known Unknown

Is 3D reconstruction unique?



The solution is defined (at least) up to a global **projective transformation**:

$$m_{ij} \simeq P_i M_j = P_i \underbrace{QQ^{-1}}_{identity} M_j = \underbrace{P_i Q}_{identity} \underbrace{Q^{-1} M_j}_{identity}$$

If cameras are calibrated, then the reconstruction ambiguity is represented (at least) by a global **rotation, translation and scale.**

The task of solvability is to analyse the **ambiguities** inherent to the SfM problem:

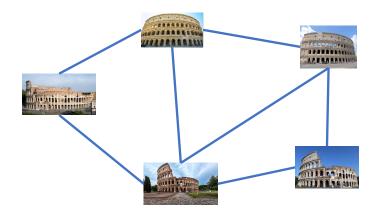
- single transformation \rightarrow well-posed problem \checkmark
- multiple transformations \rightarrow ill-posed problem X

There are many ways to approach SfM!

Here we focus on a framework that recovers **camera motion** from two-view relationships only (no points):

- Essential matrix (calibrated)
- Fundamental matrix (uncalibrated)

The problem can be represented as a **viewing graph**:

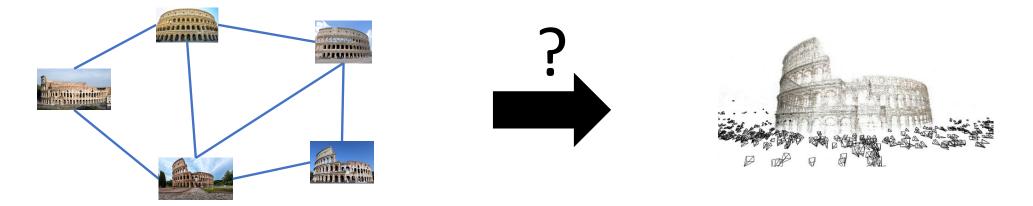


- Nodes = cameras/images
- Edges = two-view relations

Levi & Werman. *The viewing graph*. CVPR 2003.



For which graphs do we have a **well-posed** problem?



A graph is called **solvable** if and only if the available two-view relationships **uniquely** (up to a single transformation) determine the cameras \rightarrow *unique solution* **X** Otherwise it is called **non solvable** \rightarrow *multiple (infinitely many) solutions*

Here we focus on **solvability** only (*we do not address reconstruction*).

	Calibrated	Uncalibrated	
Solvability	Arrigoni & Fusiello. Bearing-based network localizability: a unifying view. IEEE TPAMI (2019).	 Levi & Werman. The viewing graph. CVPR 2003. Rudi, Pizzoli & Pirri. Linear solvability in the viewing graph. ACCV 2011. Trager, Osserman, & Ponce. On the solvability of viewing graphs. ECCV 2018. Arrigoni, Fusiello, Ricci & Pajdla. Viewing graph solvability via cycle consistency. ICCV (2021). 	
Reconstruction	Ozyesil, Voroninski, Basri & Singer. A survey of structure from motion. Acta Numerica (2017).	Kasten, Geifman, Galun & Basri. GPSfM: global projective SfM using algebraic constraints on multi-view fundamental matrices. CVPR (2019)	

It is important to check solvability **before running SfM:**

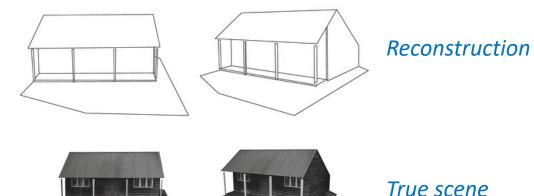
If the graph is solvable, the SfM problem is well-posed.

× If the graph is not solvable, the problem is ill-posed: no method will return a useful solution.

Outline

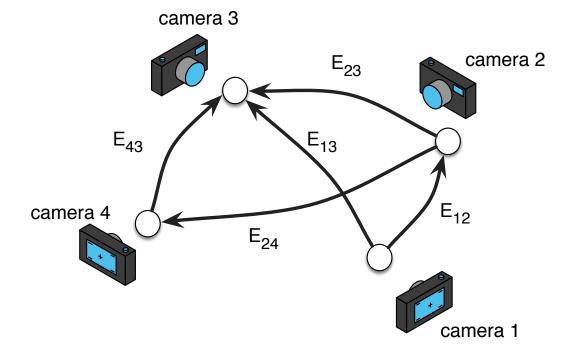
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- Calibration matrix is required in advance
- Reconstruction is **metric** (up to scale)



The Calibrated Case Problem Formulation

The **viewing graph** is a graph where vertices correspond to cameras and edges represent essential matrices.



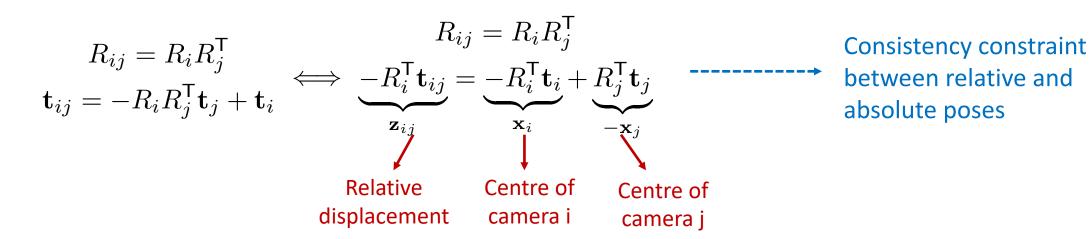
Each essential matrix can be decomposed into:

- Relative rotation R_{ii}
- *Relative translation* t_{ij} (known up to scale)

The Calibrated Case Problem Formulation

Solvable graph \Leftrightarrow two-view transformations uniquely (up to a *single* rotation, translation & scale) determine the camera poses.

- We consider a **noiseless**-case
- We split the problem into rotation and translation:



The Calibrated Case Problem Formulation

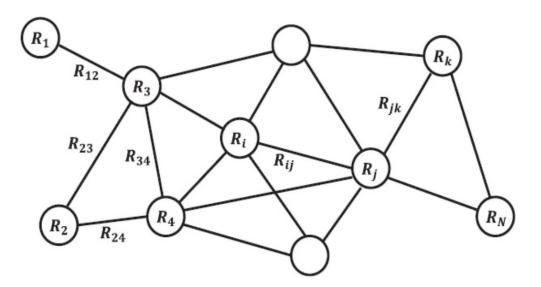
Solvable graph \Leftrightarrow two-view transformations uniquely (up to a *single* rotation, translation & scale) determine the camera poses.

- We consider a **noiseless**-case
- We split the problem into rotation and translation:

 \leftarrow The magnitude of relative translations are unknown: $||\mathbf{t}_{ij}|| = ||\mathbf{z}_{ij}|| = ?$

The Calibrated Case Rotations

In which cases can we uniquely (up to a global rotation) recover camera rotations starting from relative rotations?

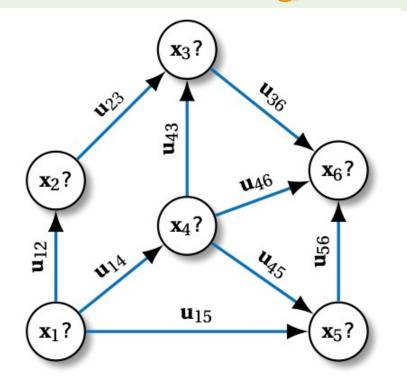


Given a **spanning tree**, a solution can be found by setting the root to the identity and propagating the consistency constraint:

$$R_i = R_{ij}R_j \iff R_{ij} = R_i R_j^T$$

Solvability for rotations ⇔ **connected** viewing graph

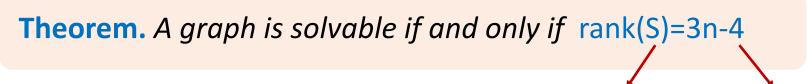
In which cases can we uniquely (up to translation & scale) recover camera positions from pairwise directions?



• **Nodes** = unknown locations

$$\mathbf{u}_{ij} = \frac{\mathbf{x}_i - \mathbf{x}_j}{||\mathbf{x}_i - \mathbf{x}_j||} \iff \mathbf{u}_{ij} \times (\mathbf{x}_i - \mathbf{x}_j) = 0$$

A solution can be found from the direction constraint, which is a **linear** equation!



Localization Translation & Equation: Sx=0 scale ambiguity

If the viewing graph is **solvable**, then the problem is well-posed.

 \times Otherwise, the problem is ill-posed: the **largest solvable component** has to be extracted \Leftrightarrow clustering rows in the null-space of S

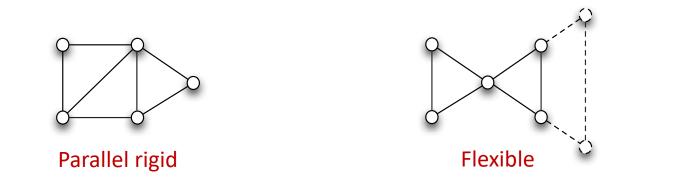
F. Arrigoni, A. Fusiello. *Bearing-based network localizability: a unifying view.* IEEE TPAMI (2019).

W. Whiteley. *Matroids from Discrete Geometry*. American Mathematical Society (1997)

R. Kennedy, K. Daniilidis, O. Naroditsky, C. J. Taylor. Identifying maximal rigid components in bearing-based localization. IROS (2012)

Solvability for translations ⇔ **parallel rigid** viewing graph

Definition. A graph is **parallel rigid** when all the configurations with parallel edges differ by translation and scale. Otherwise it is called **flexible**.

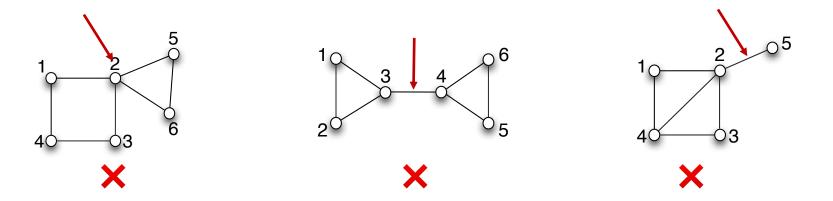


This is a well studied task!

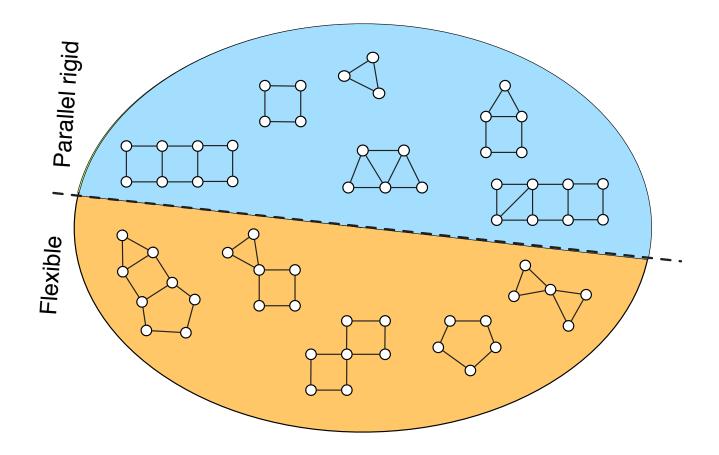
I O. Ozyesil, A. Singer. Robust camera location estimation by convex programming. CVPR (2015).

A parallel-rigid graph must satisfy the following **necessary conditions**:

- it has at least (3n-4)/2 edges
- It is **bridgeless** (i.e., it remains connected after removing any edge).
- It is **biconnected** (i.e. it does not have **articulation points** meaning that it remains connected after removing any node).



The Calibrated Case Examples



- A single cycle of length 3 or 4 is parallel rigid, whereas longer cycles are flexible
- Union of rigid graphs with a common edge is also rigid ⇒ sufficient conditions

The Calibrated Case Examples

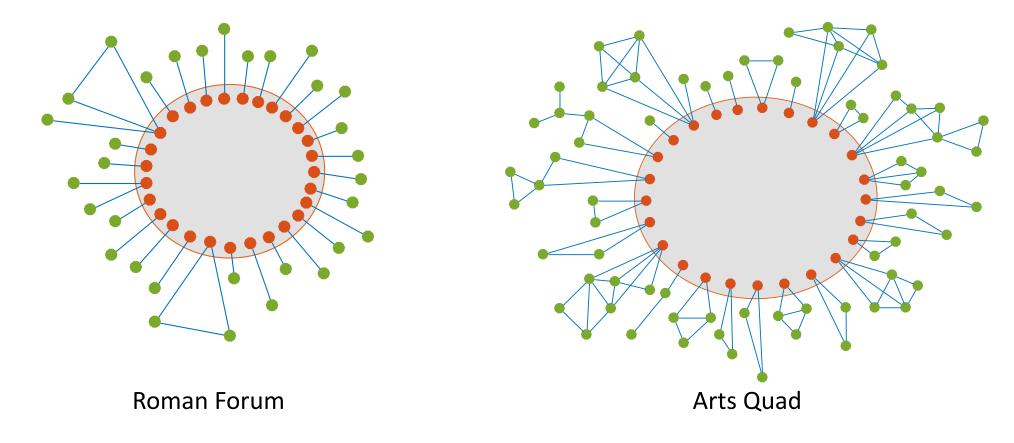
nodes	% edges	rigid	articulation	bridges
5530	2	×	30	10
2508	10	×	59	62
1134	11	×	28	28
930	6	×	60	68
918	25	×	19	20
627	50	×	17	19
553	68	\checkmark	_	_
508	19	×	19	19
474	47	×	7	7
458	26	×	9	10
394	31	×	17	15
376	29	×	17	18
354	40	×	8	9
247	67	×	6	7
	5530 2508 1134 930 918 627 553 508 474 458 394 376 354	5530 2 2508 10 1134 11 930 6 918 25 627 50 553 68 508 19 474 47 458 26 394 31 376 29 354 40	5530 2 X 2508 10 X 1134 11 X 930 6 X 918 25 X 627 50 X 553 68 ✓ 508 19 X 474 47 X 458 26 X 394 31 X 376 29 X 354 40 X	5530 2 X 30 2508 10 X 59 1134 11 X 28 930 6 X 60 918 25 X 19 627 50 X 17 553 68 \checkmark $ 508$ 19 X 19 474 47 X 7 458 26 X 9 394 31 X 17 376 29 X 17 354 40 X 8

Cornell ArtsQuad http://vision.soic.indiana.edu/projects/disco/

1DSfM datasets http://www.cs.cornell.edu/projects/1dsfm/

The Calibrated Case Examples

Simplified representation: edges outside the largest rigid component are drawn.



The Calibrated Case Summary

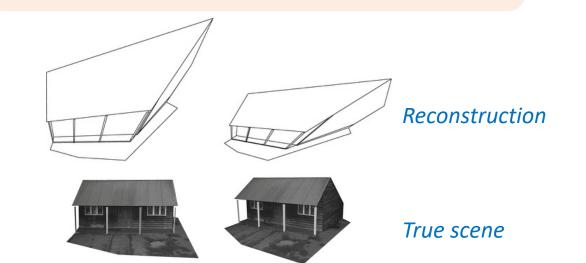
Solvability for rotations ⇔ **connected** viewing graph Solvability for translations ⇔ **parallel rigid** viewing graph

- Parallel rigidity can be tested from the rank of a linear system.
- Maximal components can be extracted from the **null-space** of such a system.
- Large-scale datasets can be processed.

Outline

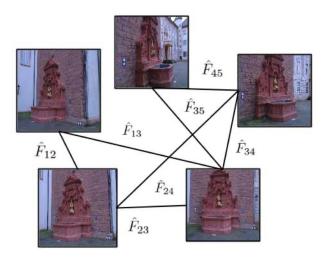
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- No assumptions
- Reconstruction is **projective**



The Uncalibrated Case Problem Formulation

The **viewing graph** is a graph where vertices correspond to cameras and edges represent fundamental matrices.

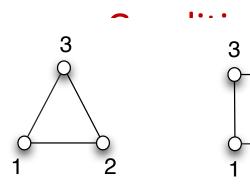


- Solvability depends on the graph and camera centres only.
- It can be reduced to a property of the graph only if we assume **generic** centres.

Solvable graph \Leftrightarrow it uniquely (up to a *single* projective transformation) determines a projective configuration of cameras.

The Uncalibrated Case

- A solvable graph has at least (1
- In a solvable graph, all the v adjacent vertices have degree



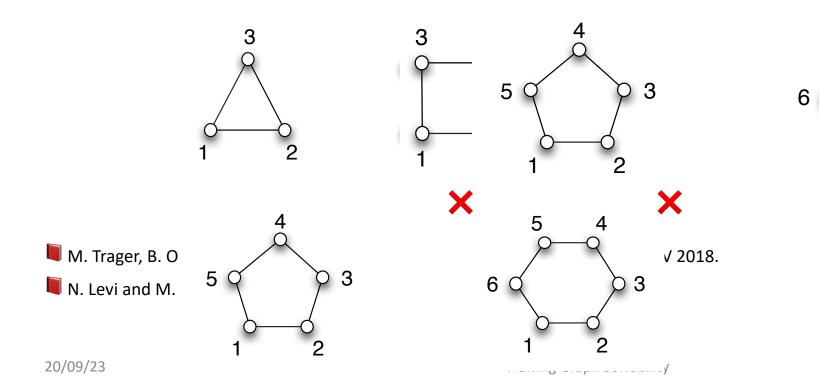
4

5

4

3

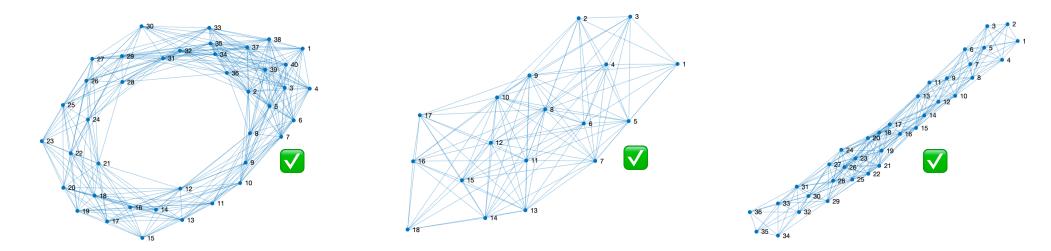
nd no two



Nec

The Uncalibrated Case Sufficient Conditions

- Triangulated graphs are solvable
- Constructive approaches are also available



M. Trager, M. Hebert, and J. Ponce. *The joint image hand-book*. ICCV 2015.

A. Rudi, M. Pizzoli, and F. Pirri. *Linear solvability in the viewing graph*. ACCV 2011.

The Uncalibrated Case Algebraic Characterization

Idea: characterize the set of projective transformations that represent all possible ambiguities of the problem.

First, let us identify the family of transformations that leave a **single camera** fixed.

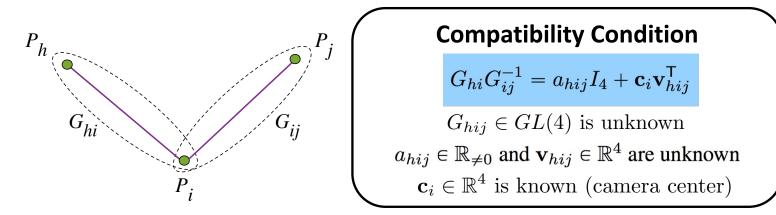
Proposition. Let P be a camera with centre c. All the solutions to PG = aPfor $G \in GL(4, \mathbb{R})$ and $a \in \mathbb{R}_{\neq 0}$ are given by $G = aI_4 + \mathbf{cv}^{\mathsf{T}} \quad \forall a \in \mathbb{R}_{\neq 0}, \mathbf{v} \in \mathbb{R}^4$

M. Trager, B. Osserman, and J. Ponce. *On the solvability of viewing graphs*. ECCV 2018.

The Uncalibrated Case Algebraic Characterization

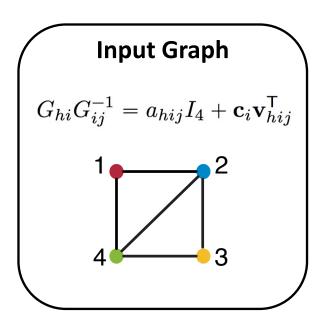
What happens when we have *multiple cameras*, represented as a viewing graph?

Let us assign an unknown projective transformation G_{ij} to every edge, and let us consider two edges (h, i) and (i, j) with a common vertex i.



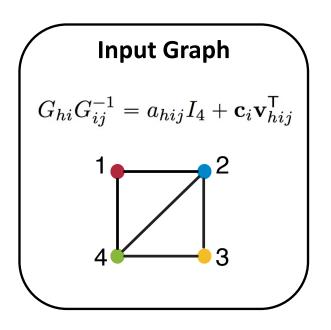
M. Trager, B. Osserman, and J. Ponce. On the solvability of viewing graphs. ECCV 2018.

The Uncalibrated Case Algebraic Characterization



Polynomial system of equations with many unknowns $G_{hi} \in GL(4)$ is unknown $a_{hij} \in \mathbb{R}_{\neq 0}$ and $\mathbf{v}_{hij} \in \mathbb{R}^4$ are unknown $\mathbf{c}_i \in \mathbb{R}^4$ is known (camera center) (h, i) and (i, j) are adjacent edges

M. Trager, B. Osserman, and J. Ponce. *On the solvability of viewing graphs*. ECCV 2018.



Polynomial system of equations with many unknowns

 $G_{hi} \in GL(4)$ is unknown

 $a_{hij} \in \mathbb{R}_{\neq 0}$ and $\mathbf{v}_{hij} \in \mathbb{R}^4$ are unknown

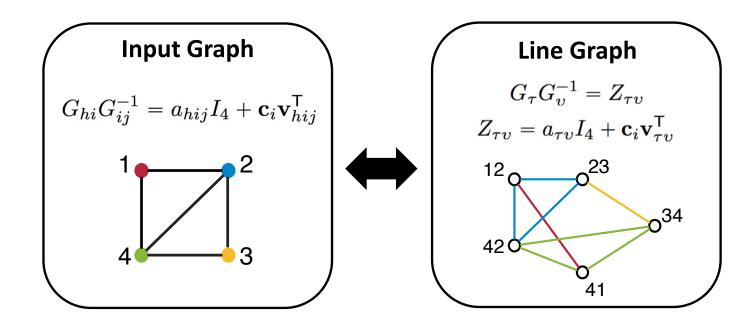
 $\mathbf{c}_i \in \mathbb{R}^4$ is known (camera center)

 $\left(h,i\right)$ and $\left(i,j\right)$ are adjacent edges

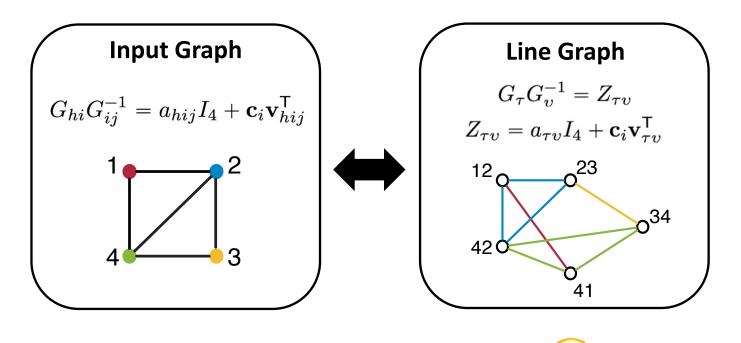
M. Trager, B. Osserman, and J. Ponce. *On the solvability of viewing graphs*. ECCV 2018.

• It is possible eliminate variables 😂

Arrigoni, Fusiello, Ricci & Pajdla. Viewing graph solvability via cycle consistency. ICCV (2021).



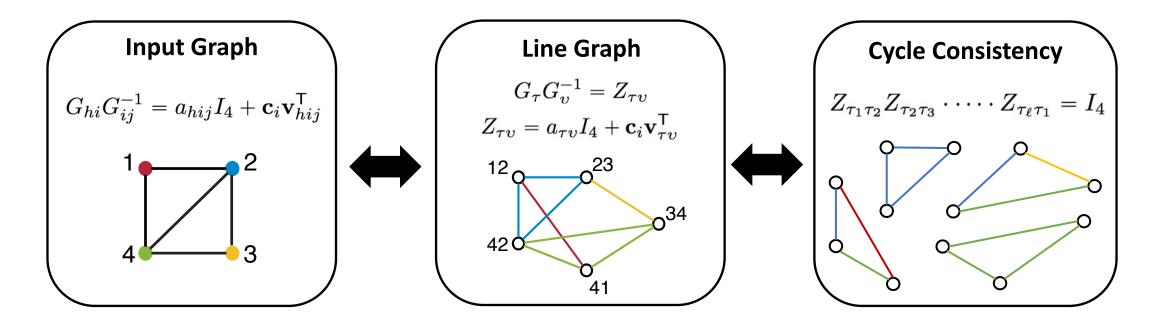
- Each node is an edge in the input graph;
- Two nodes are linked if the corresponding edges are adjacent in the input graph.
- There is one equation for each edge in the line graph.



- Each node is an edge in the input graph;
- Two nodes are linked if the corresponding edges are adjacent in the input graph.
- There is one equation for each edge in the line graph.

How can we eliminate the G variables?

Idea:
$$Z_{12,23} \cdot Z_{23,42} \cdot Z_{42,12} = G_{12} \underbrace{G_{23}^{-1}G_{23}}_{I} \underbrace{G_{42}^{-1}G_{42}}_{I} G_{12}^{-1} = I$$



cycle consistency (on all cycles) \Leftrightarrow cycle consistency (on a basis)

The Uncalibrated Case Algorithm

Algorithm 1 Viewing Graph Solvability

Input: undirected graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$

Output: solvable or not solvable

- 1. randomly sample the camera centres
- 2. compute the line graph $\mathcal{L}(\mathcal{G})$
- 3. compute a cycle consistency basis for $\mathcal{L}(\mathcal{G})$
- 4. set up equations

5. compute the number s of real solutions -----

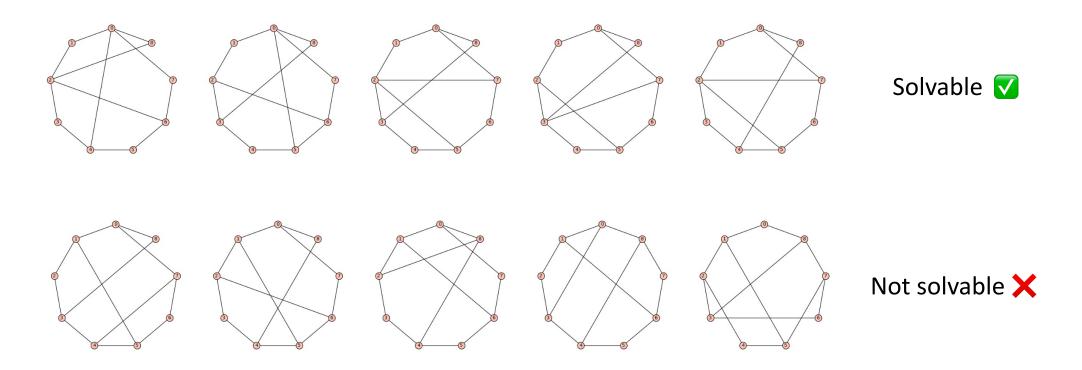
if s = 1 then

solvable else not solvable end if **Gröbner basis** (symbolic computation)

https://github.com/federica-arrigoni/solvability

The Uncalibrated Case Examples

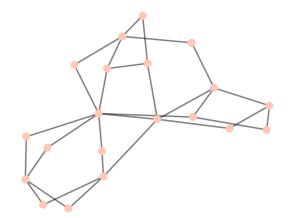
Minimal viewing graphs with 9 vertices

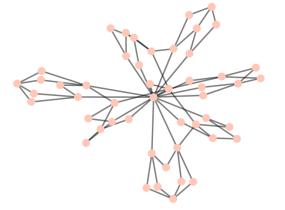


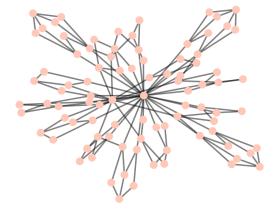
The Uncalibrated Case Examples

Execution times on minimal graphs

Nodes	10	20	30	40	50	60	70	80	90
Time	1.6 s	9 s	93 s	3 min	15 min	35 min	1 h	$\approx 2 h$	>4 h







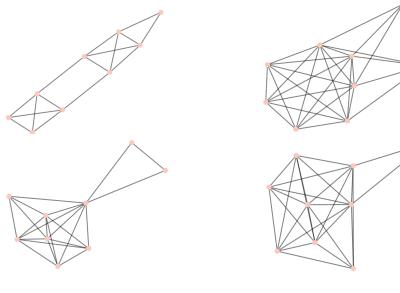
Solvable graph with 90 nodes

Solvable graph with 20 nodes

Solvable graph with 50 nodes

The Uncalibrated Case Examples

Subgraphs with 9 nodes sampled from real structure-from-motion viewgraphs



Unsolvable

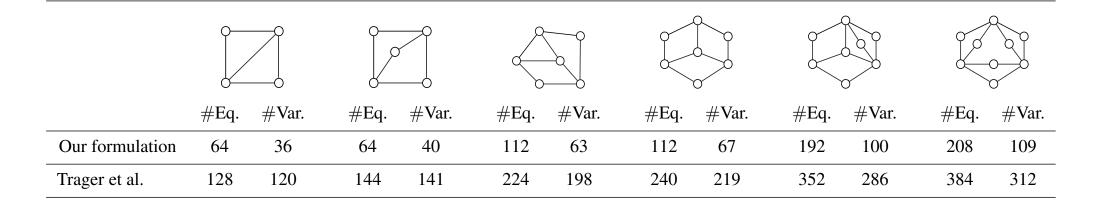


Solvable

		Solvable	Unsolvable			
Data set	by suff.	by Alg. <mark>1</mark>	Tot.	by nec.	by Alg. <mark>1</mark>	Tot
Alcatraz Courtyard	200	0	200	0	0	0
Buddah Tooth	178	20	198	2	0	2
Pumpkin	169	22	191	8	1	9
Skansen Kronan	179	8	187	13	0	13
Tsar Nikolai I	196	0	196	4	0	4
Alamo	136	16	152	48	0	48
Ellis Island	136	30	166	34	0	34
Gendarmenmarkt	128	11	139	61	0	61
Madrid Metropolis	88	28	116	84	0	84
Montreal Notre Dame	140	12	152	48	0	48
Notre Dame	165	18	183	17	0	17
NYC Library	110	19	129	71	0	71
Piazza del Popolo	105	22	127	73	0	73
Piccadilly	109	23	132	68	0	68
Roman Forum	114	28	142	58	0	58
Tower of London	123	18	141	59	0	59
Trafalgar	86	16	102	98	0	98
Union Square	74	19	93	107	0	107
Vienna Cathedral	122	8	130	70	0	70
Yorkminster	116	14	130	70	0	70
Cornell Arts Quad	76	23	99	101	0	101

The Uncalibrated Case Summary

• Thanks to cycle consistency, **less unknowns** are involved than previous work:

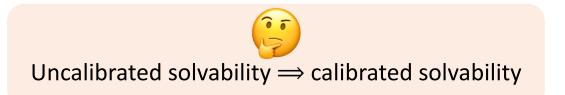


- It is possible to classify previously undecided viewing graphs and extend solvability testing up to minimal graphs with 90 nodes.
- Larger/denser graphs can not be processed 😕



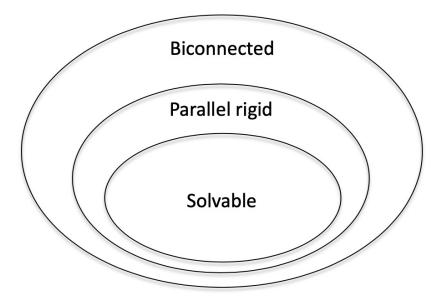
Outline

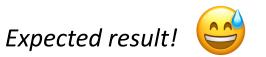
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Calibrated vs Uncalibrated

Proposition. A solvable (uncalibrated) graph is parallel rigid.





Well-posed with uncalibrated cameras \Rightarrow well-posed with calibrated cameras

Arrigoni, Fusiello, Rizzi, Ricci & Pajdla. Revisiting viewing graph solvability: an effective approach based on cycle consistency. TPAMI (2022).

Calibrated vs Uncalibrated

Proposition. A solvable (uncalibrated) graph is parallel rigid.

Proof [sketch]. Parallel rigid graph \Leftrightarrow for any partition of the edges: $\sum_{i=1}^{k} (3|\mathcal{V}_i| - 4) \ge 3n - 4$ Solvable graph \Rightarrow for any partition of the edges: $\sum_{i=1}^{k} (11|\mathcal{V}_i| - 15) \ge 11n - 15$ *Only necessary condition! Unknown if the opposite holds*

Arrigoni, Fusiello, Rizzi, Ricci & Pajdla. Revisiting viewing graph solvability: an effective approach based on cycle consistency. TPAMI (2022).

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Conclusion

	Calibrated	Uncalibrated
Formulation	Linear system	Polynomial system
Datasets	Large-scale	Small-scale
Interpretation	Connected + Parallel rigid	?
Components	Null-space computation	?
I		
	"Solved"	Open issues

References

F. Arrigoni, T. Pajdla & A. Fusiello. *Viewing graph solvability in practice*. ICCV (2023).

F. Arrigoni, A. Fusiello, R. Rizzi, E. Ricci & T. Pajdla. *Revisiting viewing graph solvability: an effective approach based on cycle consistency*. IEEE TPAMI (2022).

F. Arrigoni, A. Fusiello, E. Ricci & T. Pajdla. *Viewing graph solvability via cycle consistency*. ICCV (2021). **Best paper honourable mention Y**

F. Arrigoni & A. Fusiello. *Bearing-based network localizability: a unifying view.* IEEE TPAMI (2019).

Thank you for your attention!

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